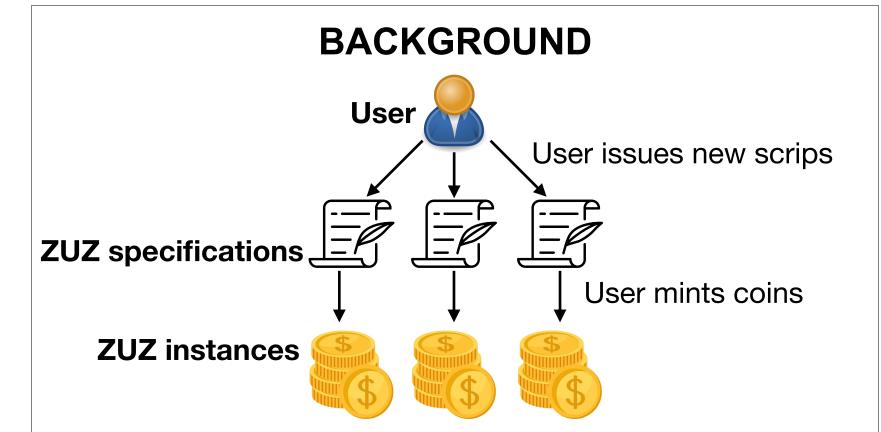
Privacy-Preserving Reputation-Based Lending Systems

ABSTRACT

Concerned with the privacy of their digital transactions, users are averse to leaving the very digital trace necessitated by reputation-based lending systems. Our design converts between account and UTXO wallet representations to hide the transfer's origin, value, and destination while the underlying scrip still gains reputation. Our approach results in fast and cheap online computation, while total transaction times scale linearly in the number of recipients.



In a *reputation-based lending system*, users can issue their own scrips, or ZUZ specifications, to be redeemed in exchange for goods and services in the future. In turn, these specifications gain value when *instances* of them are traded by individuals or at businesses outside of the ones that issued them. When such transactions are made transparent and universally accessible, such as by being recorded on a *distributed public ledger*, we can base a scrip's reputation off of its acceptability.

Under such a system, a user can conduct an anonymous transaction by minting ZUZ instances from a new specification issued under a new pseudonym. Since neither the specification nor the pseudonym have any history associated with them, the transaction is completely anonymous; however, for the very same reason, other users on the network have little to no incentive to accept these funds in a transaction. Without a way for users to spend their existing funds privately, such a system is rendered unusable.

In our smart contract implementation, users can create ZUZ specifications, and mint, *pour*, and transfer ZUZ instances. A pour operation allows a user to convert a public ZUZ instance into a private ZUZ instance, and vice versa.

Public instances are represented by numerical balances, like accounts, whereas private instances are represented as lists of commitments, like unspent transaction outputs.

Note that the ZUZ specification is made public in all parts of the interface so that the underlying scrip can still gain reputation through private transfers.

- and

by providing a zero-knowledge proof π_{POUR} which shows • the sender can unlock all the provided commitments, and • balance is preserved

<u>Transfer</u> (*spec*, $[cm_1^{recv} \dots cm_r^{recv}]$, $[cm_1^{old} \dots cm_w^{old}]$, cm^{new} , $[Enc_{pk_1}(\{b_1, \omega_1\}) \dots Enc_{pk_r}(\{b_n, \omega_r\})]$, $\pi_{TRANSFER}$): For the ZUZ specification *spec*, the sender

- adds private balance cm_i^{recv} to recipient pk_i along with the parameters to unlock the commitment $Enc_{pk_i}(\{b_i, \omega_i\})$
- replaces sender's private balances $[cm_1^{old} \dots cm_w^{old}]$ with new private balance cm^{new}
- by providing a zero-knowledge proof $\pi_{TRANSFER}$ which shows • the sender has sufficient funds, and
- balance is preserved

In both cases, the smart contract must validate that the inputs to the ZKP match the ledger state to prevent double-spending attacks.

One remaining challenge is mitigating *wash* attacks on the reputation of ZUZ specifications under private transfers.

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INTERFACE

Pour (spec, b_{pub}^{old} , $[cm_1^{old} \dots cm_w^{old}]$, b_{pub}^{new} , cm^{new} , π_{POUR}): For the ZUZ specification *spec*, the sender converts between their • old public balance b_{nub}^{old} and private balances $[cm_1^{old}...cm_w^{old}]$,

• new public balance b_{nub}^{new} and private balance cm^{new}

TRANSFER PROTOCOLS One-step transfer: Smart Alice contract $\pi_{TRANSFER}$ ---- </> $cm_B, Enc_B(b)$ CM_A $c, E_{nc}(0)$ Destination-anonymous Smart contract validation • **Issue:** Alice pays for Bob's privacy <u>Two-step transfer:</u> Step 1: Bob requests Alice, Alice validates Bob Charlie Adam $cm_{A'}, Enc_A(b_{A'}, \omega_{A'}) \quad cm_B, Enc_A(b_B, \omega_B) \quad cm_C, Enc_A(b_C, \omega_C)$ 00 $cm^{D}_{TRANSFER}, \pi^{D}_{TRANSFER}$ Step 2: Alice receives funds on transfer $\pi^{A}_{TRANSFER}$ -.._ </> $cm_{TRANSFER}^{P}, cm_{A}$ Destination-anonymous • Smart contract validation of Alice and clientside validation of Bob • Bob pays for his own privacy Three-step transfer: Step 1: Alice *initiates, Bob* privacy, Alice validates validates Enc_A(b) $Enc_B(0)$ Alice $cm_B, Enc_B(b)$ 00 $\pi^{A}_{TRANSFER}$ 00 $cm_{C}, Enc_{B}(0)$ Bob Charlie Step 3: Alice receives funds on transfer • Destination- and origin-anonymous • Client-side validation of Alice and Bob

• Bob pays for his own privacy

